Optimal Personalized Filtering Against Spear-Phishing Attacks

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Malicious E-Mails



Spam

- non-targeted
- usually just a nuisance (but can waste a lot of time and money in high volumes)



Spear-phishing

- targeted
- potentially very high losses (even from a single attack)

Spear-Phishing Examples

- In 2014, a German steel mill suffered "massive" physical damage due to a cyber-attack
 - first step of the attack was spear-phishing

http://www.wired.com/2015/01/german-steel-millhack-destruction/

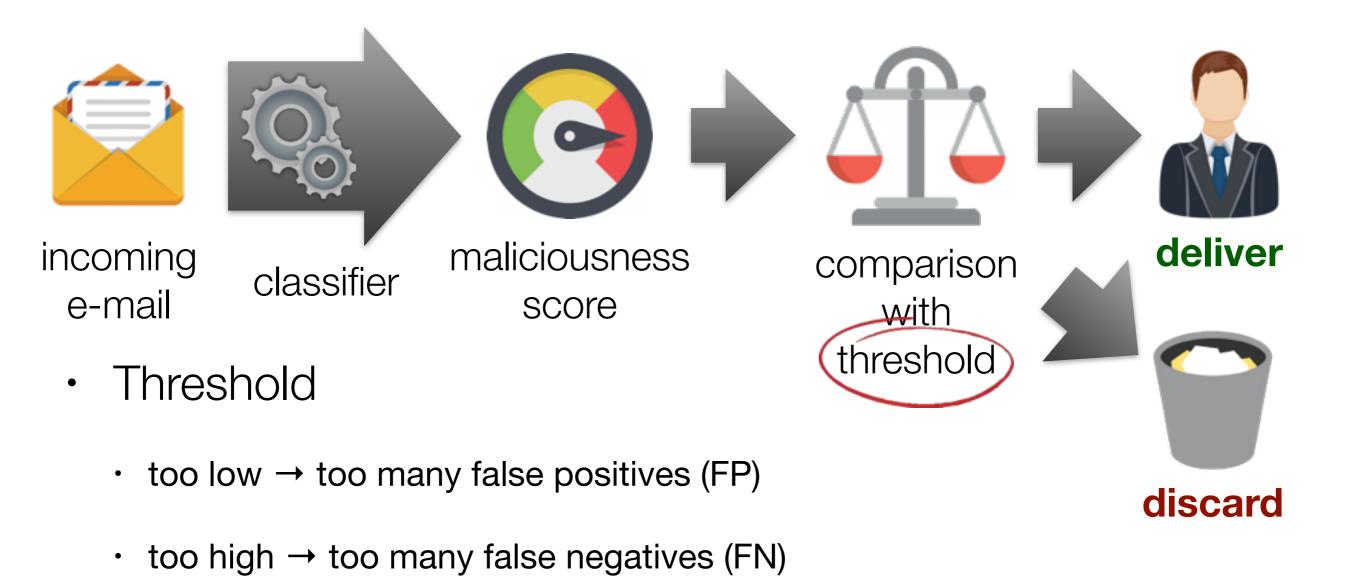
- In 2013, millions of credit and debit card accounts were compromised due to an attack against Target
 - first step of the attack was spear-phishing

http://www.huffingtonpost.com/2014/02/12/ target-hack_n_4775640.html



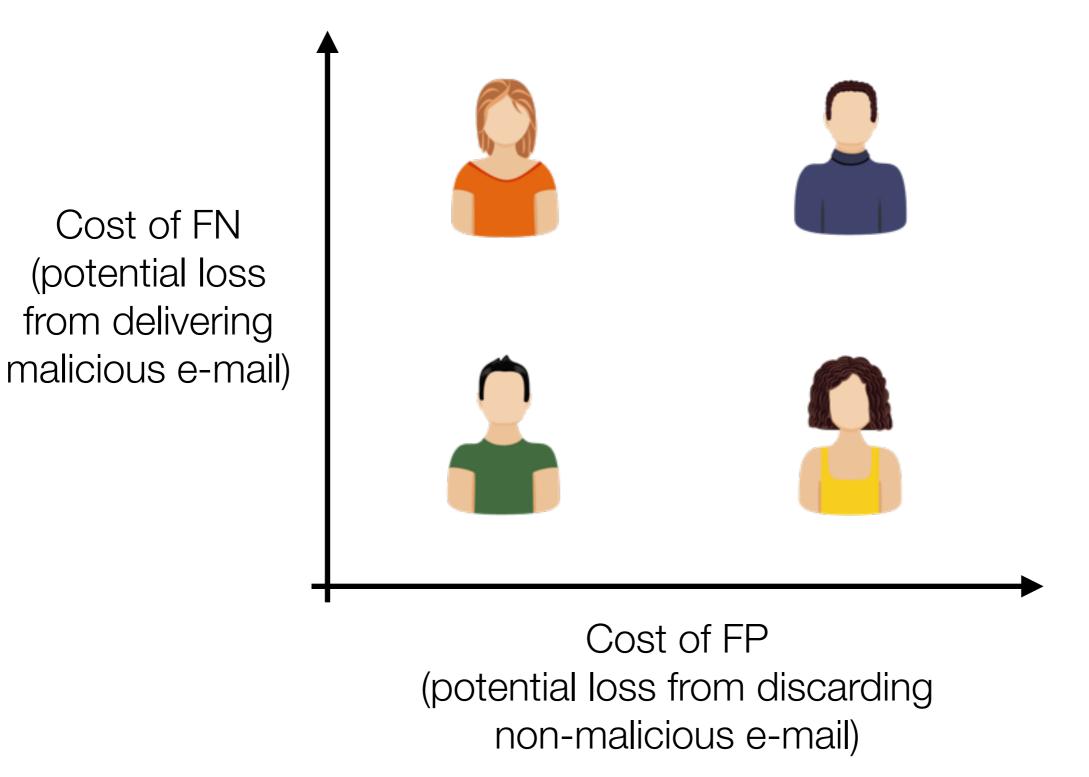


Filtering Malicious E-Mails

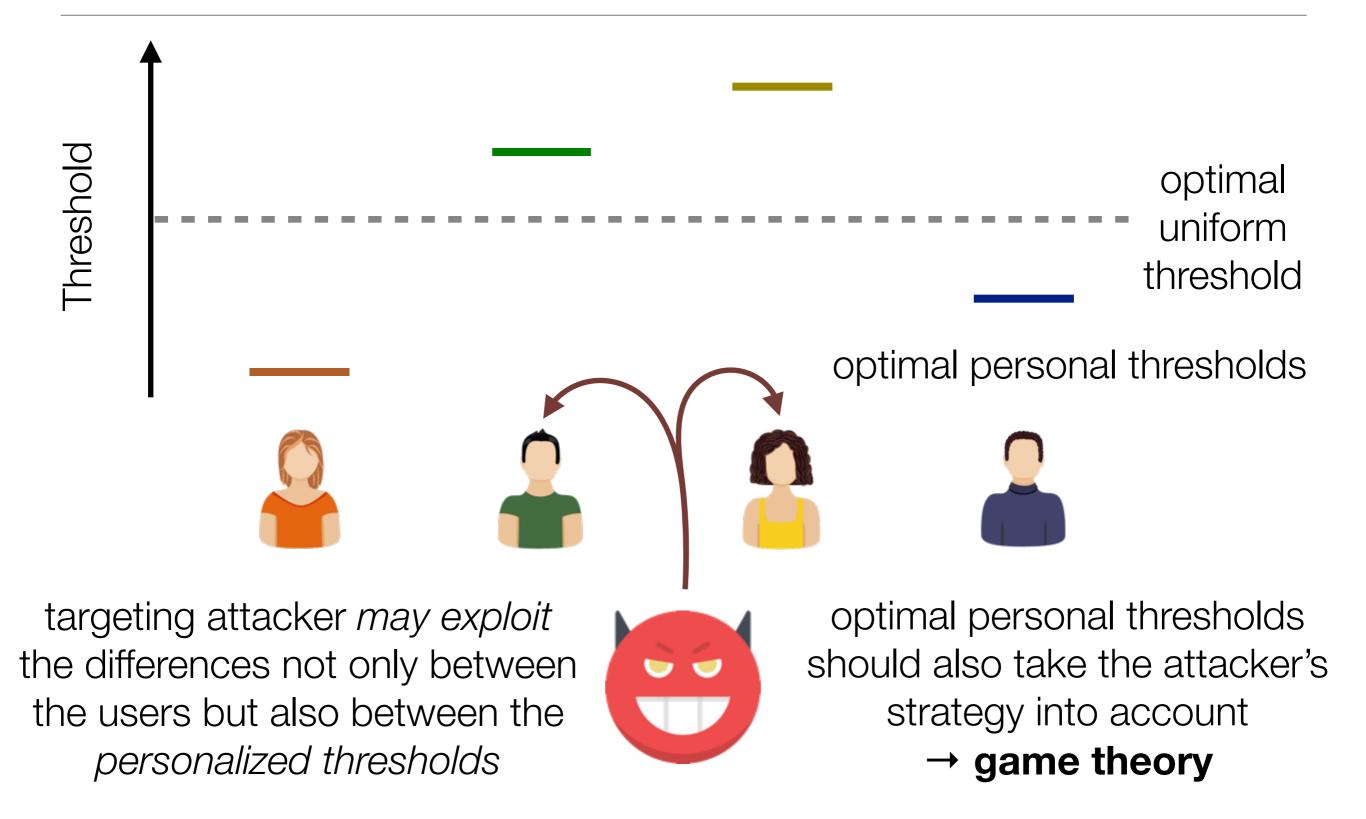


optimal value:
minimizes FP rate × cost of FP + FN rate × cost FN

Multiple Users



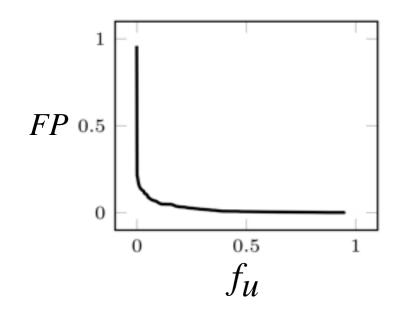
Personalized Thresholds



Game-Theoretic Model



- for each user u, selects a false negative rate f_u
- we assume that the feasible FP / FN rate pairs are given by a function $FP(f_u)$





Targeting attacker

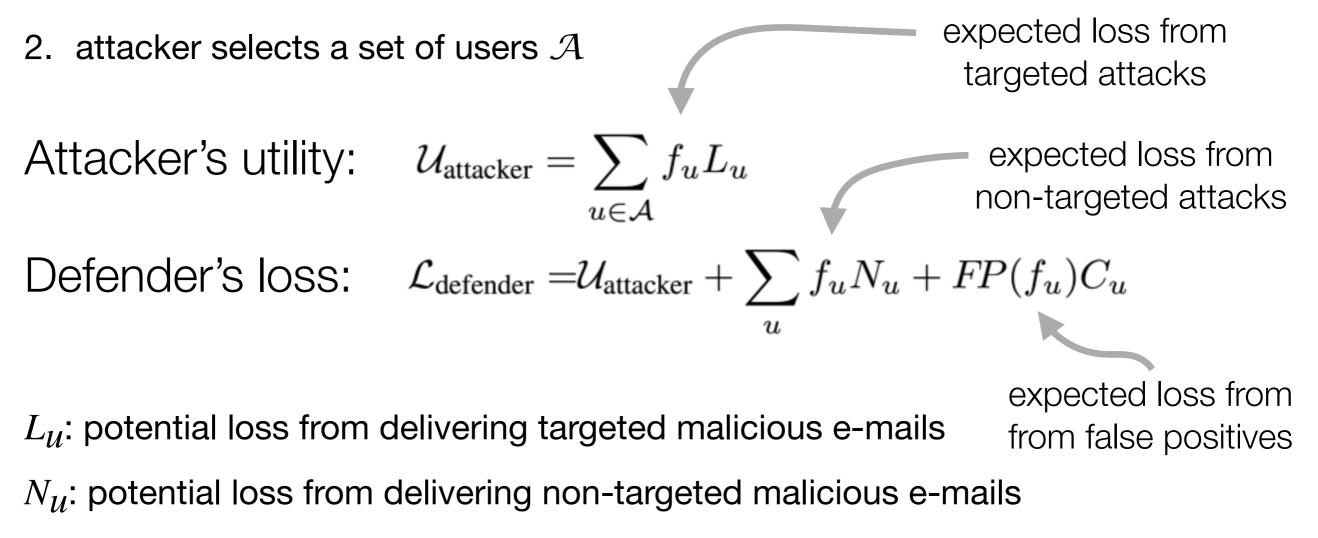
- selects a set of users A, and sends them targeted malicious e-mails
- can select at most A users (otherwise the attack is easily detected)



Game-Theoretic Model (contd.)

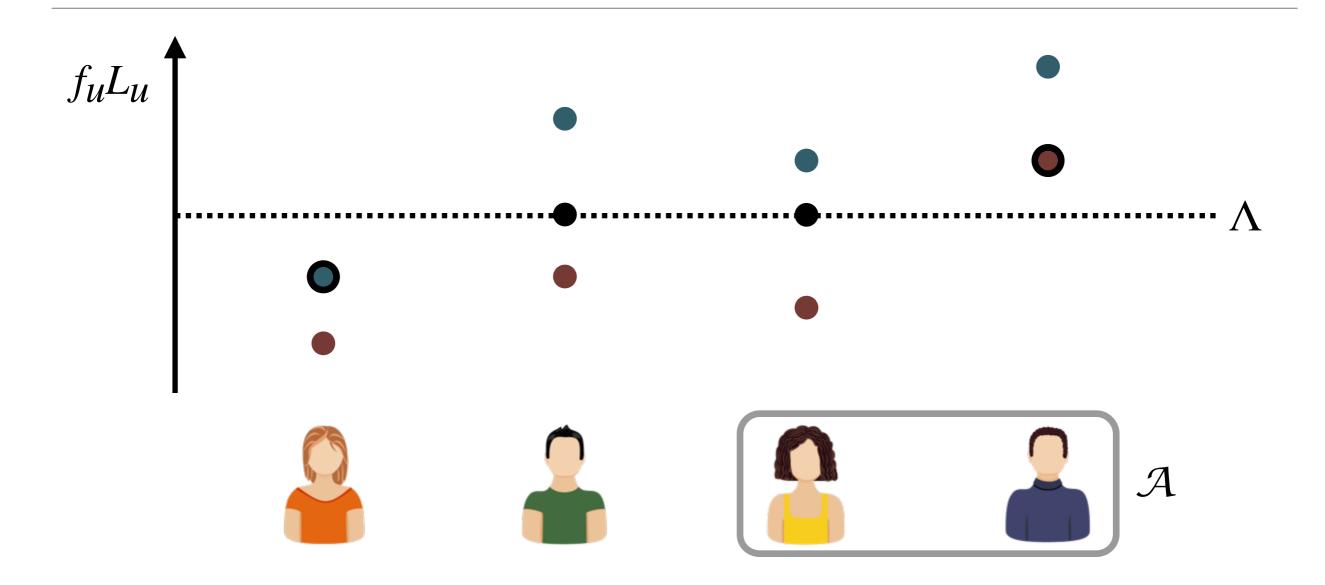
Stackelberg (leader-follower) game

1. defender selects a false negative rate f_u for each user u



 $C_{\mathcal{U}}$: potential loss from discarding non-malicious e-mails

Characterizing Optimal Strategies



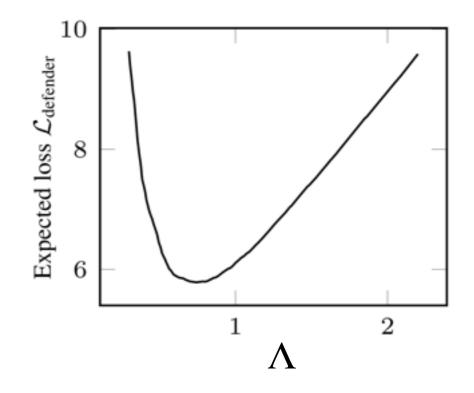
optimal value for a user given that it is <u>not selected</u> by the attacker

optimal value for a user given that it is <u>selected</u> by the attacker

Finding an Optimal Strategy

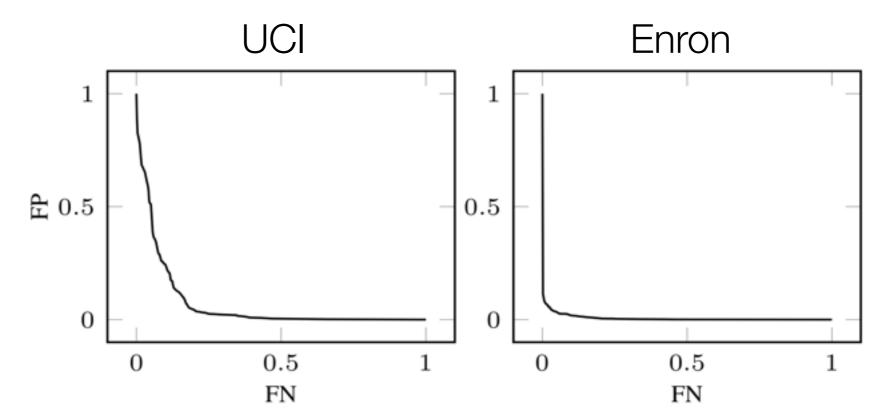
- For a given value of Λ, we can find an optimal strategy using the following polynomial-time algorithm
 - 1. For each user u, compute the loss of user u when it is not targeted as follows: if $f_u^N L_u < \Lambda$, then the loss is $f_u^N N_u + FP(f_u^N)C_u$; otherwise, the loss is $\frac{\Lambda}{L_u}N_u + FP(\frac{\Lambda}{L_u})C_u$.
 - 2. For each user u, compute the loss of user u when it is targeted as follows: if $f_u^T L_u > \Lambda$, then the loss is $f_u^T (L_u + N_u) + FP(f_u^T)C_u$; otherwise, the loss is $\frac{\Lambda}{L_u}(L_u + N_u) + FP(\frac{\Lambda}{L_u})C_u$.
 - 3. For each user u, let the cost of user u being targeted be the difference between the above computed loss values.
 - 4. Select a set A of A users with the lowest costs of being targeted.
 - 5. For every $u \in A$, let $f_u = f_u^T$ if $f_u^T L_u > \Lambda$, and let $f_u = \frac{\Lambda}{L_u}$ otherwise.
 - 6. For every $u \notin A$, let $f_u = f_u^N$ if $f_u^N L_u < \Lambda$, and let $f_u = \frac{\Lambda}{L_u}$ otherwise.
 - 7. Output the strategy f.

• Finally, we can find the optimal value of Λ using a simple binary search



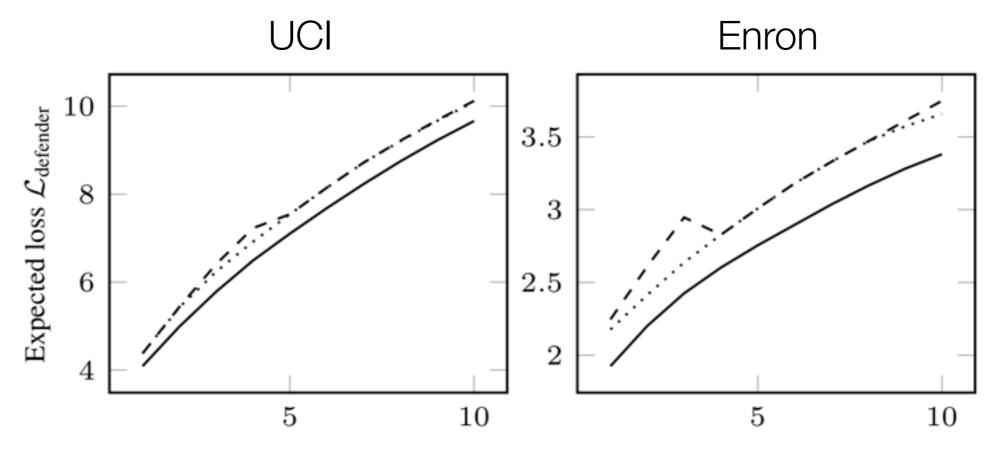
Numerical Examples

- Datasets
 - UCI Machine Learning Repository: 4601 labeled e-mails with 57 features
 - Enron dataset: 13,500 e-mails with 500 features
- Classifier: naive Bayes (note that this is just for the sake of example)
- False positive / false negative rates:



Numerical Examples - Results

• 31 users with parameter values following power-law distributions



Number of users targeted A

- optimal strategy
- --- uniform threshold not expecting strategic attacker
- ···· uniform threshold expecting strategic attacker

Conclusion & Future Work

- Conclusion
 - filtering thresholds have received less attention in the past
 - we proposed a game-theoretic model for targeted and nontargeted malicious e-mails
 - we showed how to find optimal strategies efficiently
 - numerical results show considerable improvement
- Future work
 - non-linear losses from compromising multiple users

Thank you for your attention!

Questions?

