# FlipThem: Modeling Targeted Attacks with FlipIt for Multiple Resources

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# Stealthy Attacks

- In many scenarios, attackers want to keep successful security compromises covert
- Examples

#### Cyber-espionage

 targets must not know that they are being spied on



#### Botnets

 users should not be aware that their computers are infected



# Mitigating Cover Compromises

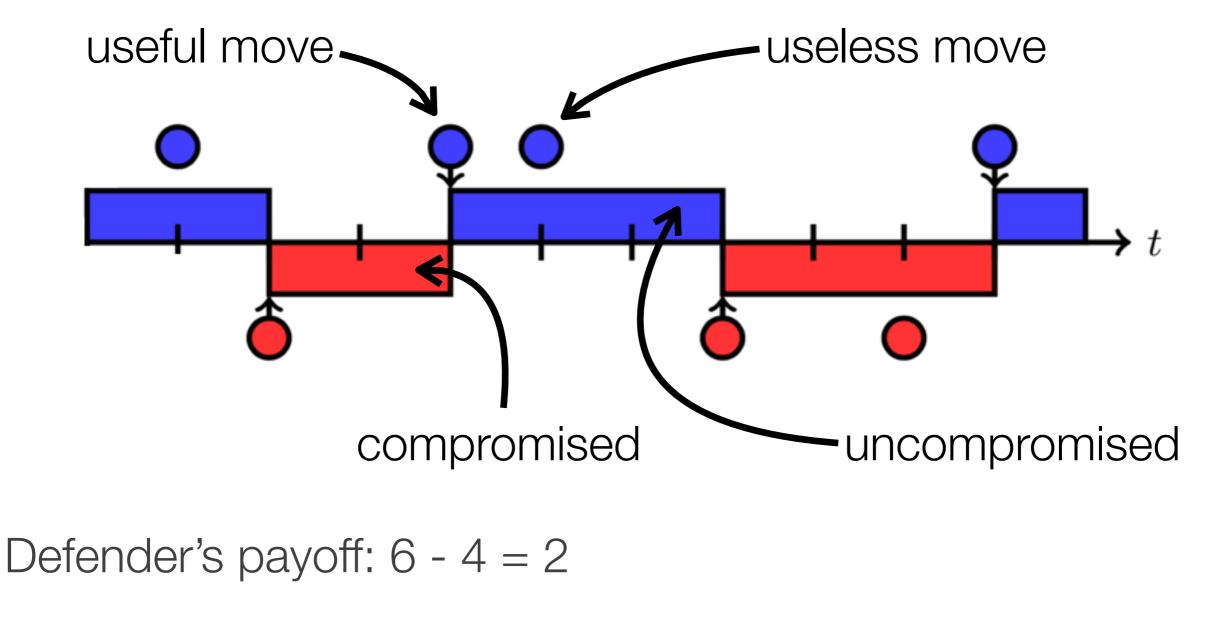
- Mitigation
  - possible losses can be minimized by resetting the computing resource into a known secure state
  - examples: changing a password or a private key, reinstalling a machine
  - "When should these moves be made?"
    - What is the optimal frequency?
    - What is the optimal scheduling?
- In practice: usually periodic key and password renewal strategies



# The FlipIt Game

- Introduced by researchers at RSA for modeling stealthy attacks against computing resources
- Resource: user account, private key, machine, etc.
- Players
  - **defender**: the rightful owner of the resource
  - **attacker**: an adversary who is trying to take over the resource
- Strategy
  - schedule for a series of costly moves (e.g., periodic)
  - each move takes control of the resource (if it is not already controlled)
- Payoff: amount of time the resource is controlled by the player - cost of moves

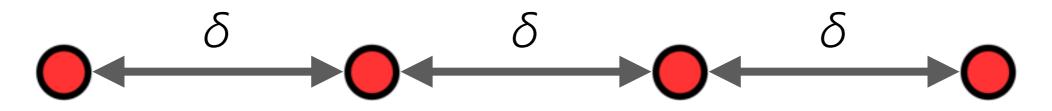
## The FlipIt Game - Graphical Illustration



Attacker's payoff: 5 - 3 = 2time controlled moves

# The FlipIt Game - Lessons Learned

• If there is no feedback, **periodic** strategies are dominant



- If the attacker learns the defender's previous moves when making a move,
  - then the defender is better off with a more random strategy, such as a renewal process with exponential interval distribution



for the attacker, periodic is still a good choice

# Multiple Resources

• FlipIt tells us how to defend a single resource



What if the security of a system depends on multiple resources?

• We could use a separate game for each resource

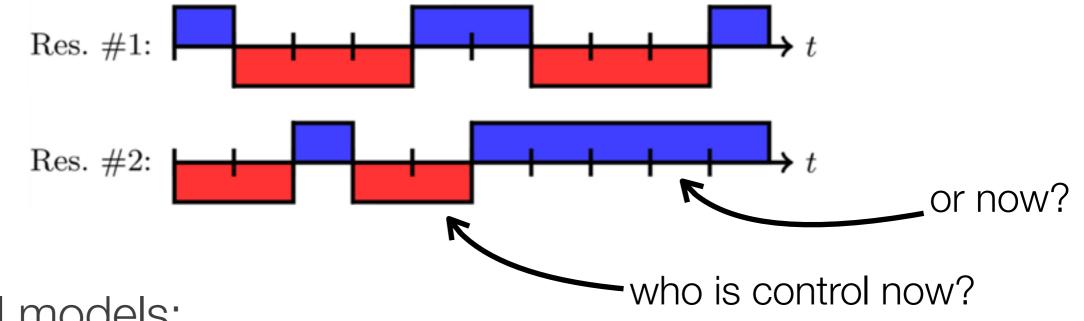


 But to exploit the dependencies between these resources, we need to model them together



# Defining the Multiple-Resource Game

- · Defining the players, the moves, etc. is straightforward
- Defining the payoffs is not straightforward



Control models:

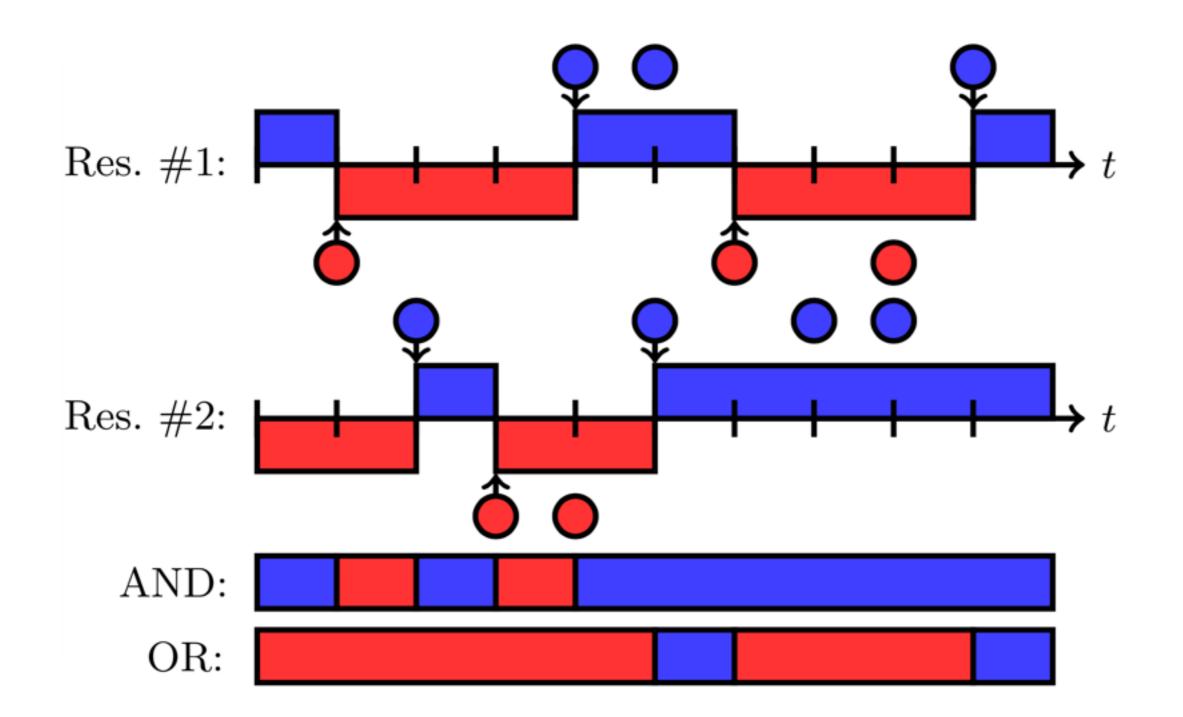
#### AND

attacker controls the system only if it controls all resources

## OR

attacker controls the system if it controls at least one resource

## Illustration of Control Models



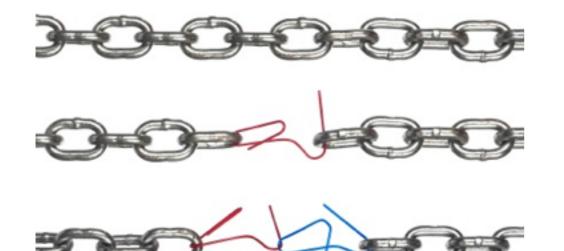
# **Control Models - Further Discussion**

## AND

- similar to the total effort model in security economics
- example: there are multiple private keys (stored separately), and the attacker needs to forge signatures for all of them
- defender is at advantage

# OR

- similar to the weakest link model in security economics
- example: there are multiple administrator accounts on a machine, and the attacker needs to compromise only one
- attacker is at advantage





# Combining Single-Resource Strategies

- Idea: build multiple-resource strategies from singleresource strategies that perform well in the FlipIt game
- Combinations:

#### Independent

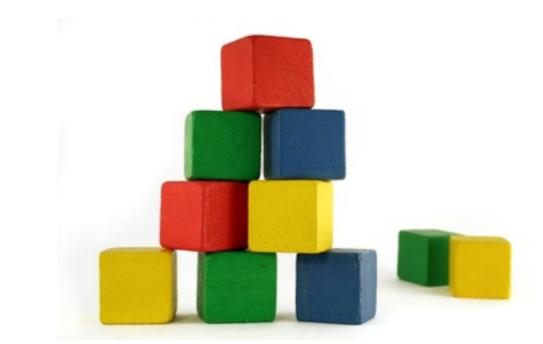
 flip each resource independently of the others (i.e., use N independent single-resource strategies)

"Which one is better?"

- For which player?
- In which control model?

## Synchronized

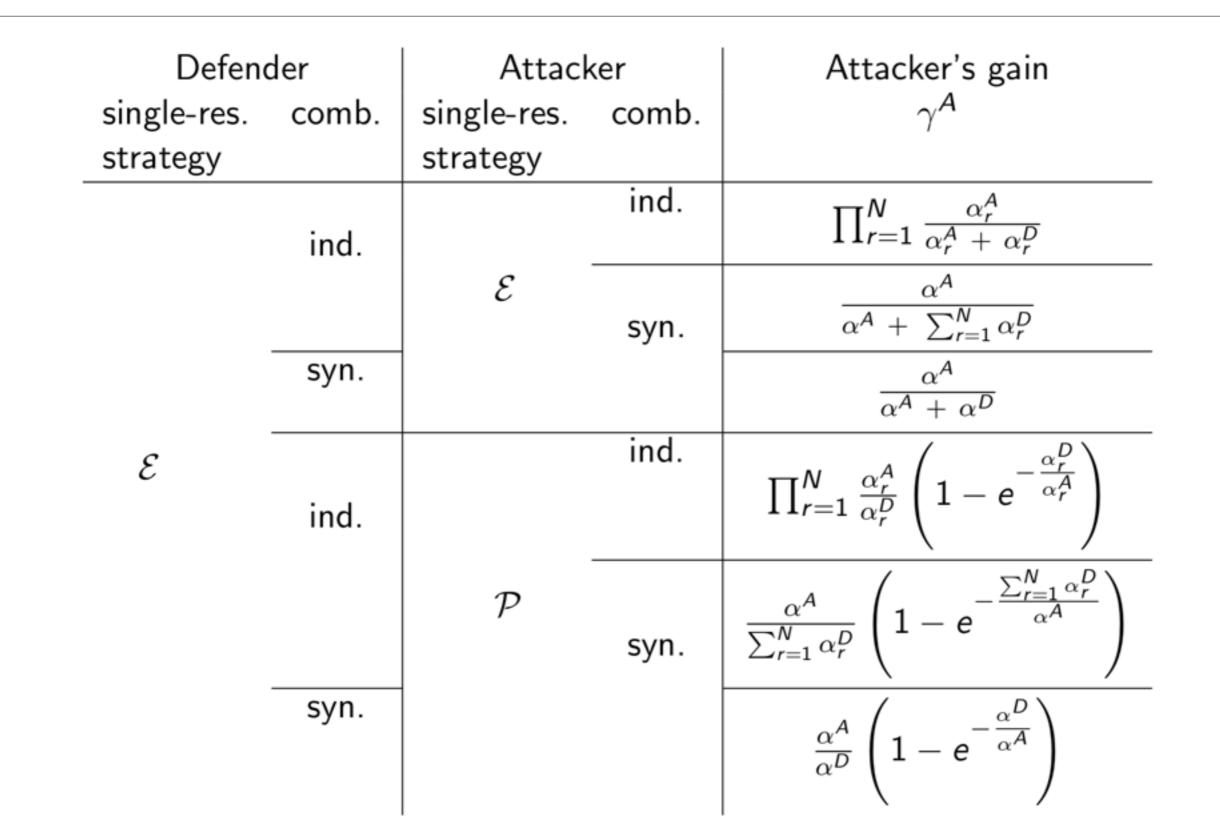
 always flip all resources together (i.e., use only one single-resource strategy for all the resources)



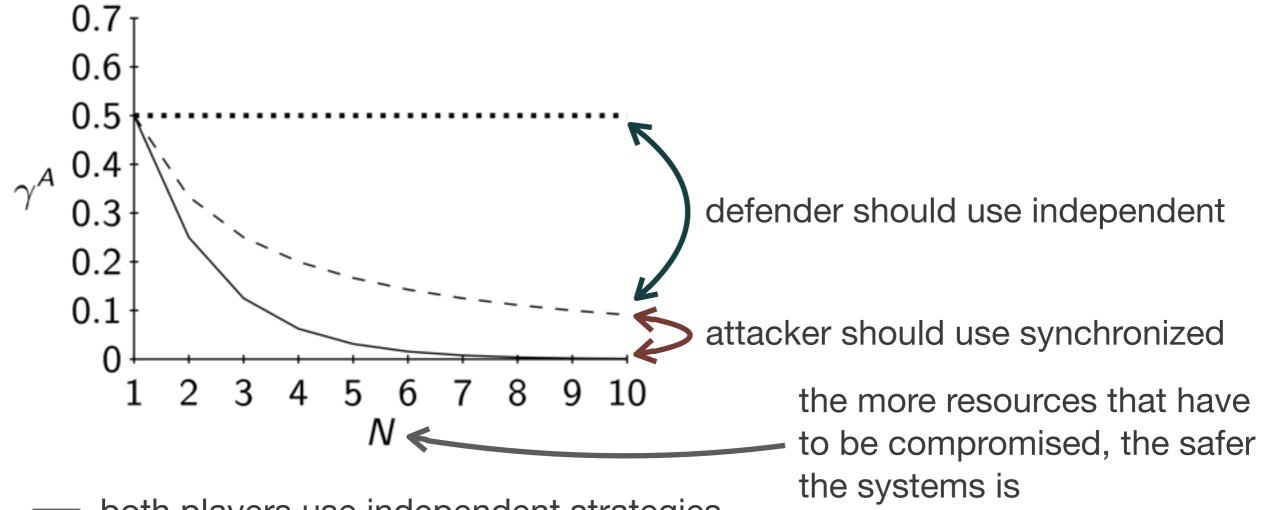
# Attacker's Gain in the AND Model - Formulae #1

Defender's combination	Attacker's combination	Attacker's gain $\gamma^{A}$
independent	independent	$\prod_{r=1}^N \int_0^\infty f_{Z^D_r}(z_r) F_{Z^A_r}(z_r) dz_r$
	synchronized	$\int_0^\infty \prod_{r=1}^N \left(1 - F_{Z^D_r}(z)\right) f_{Z^A}(z) dz$
synchronized		$\int_0^\infty f_{Z^D}(z)F_{Z^A}(z)dz$
	independent	$\int_0^\infty \prod_{r=1}^N F_{Z_r^A}(z) f_{Z^D}(z) dz$

## Attacker's Gain in the AND Model - Formulae #2



## Attacker's Gain in the AND Model - Numerical #1

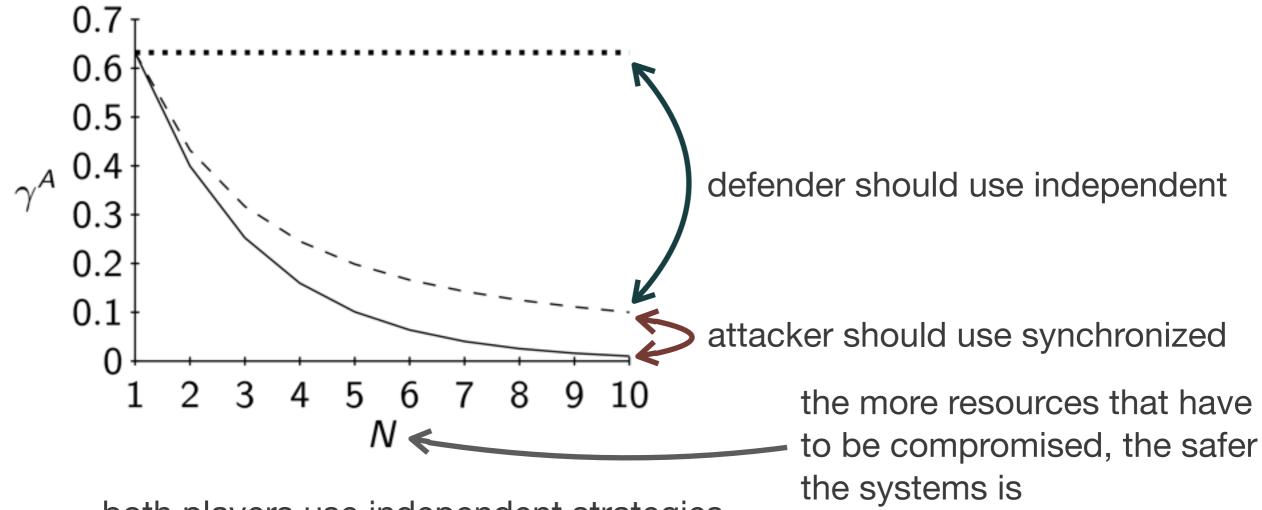


both players use independent strategies

- -- attacker uses synchronized, while defender uses independent
- •••• both players use synchronized

(both players build on exponential single-resource strategies)

## Attacker's Gain in the AND Model - Numerical #2



both players use independent strategies

- -- attacker uses synchronized, while defender uses independent
- •••• both players use synchronized

(defender builds on exponential, attacker builds on periodic single-resource strategies)

# Strategy Combinations - Lessons Learned

- In the **AND** model,
  - defender should use independent strategies
  - attacker should use synchronized strategies

Since the two control models are the same with the roles of the players reversed, we readily have that

- in the **OR** model,
  - defender should use synchronized strategies
  - attacker should use independent strategies

Modeling assumptions matter a lot!

# Markov Strategy Class

• Definition:

at each time instance, the defender may flip any subset of the resources, and the probability of flipping a given subset depends on the times elapsed since flipping each resource

- "Multi-dimensional renewal process"
- Generalizes the above single-resource combinations
  - independent: probability of flipping a given resource depends on the time elapsed since last flipping that resource, and the probability of flipping a subset is simply the product of its elements' probabilities
  - synchronized: either all resources are flipped or none are, and the probability depends on the time elapsed since the last flip

# Markov Strategies - Linear Programming Solution

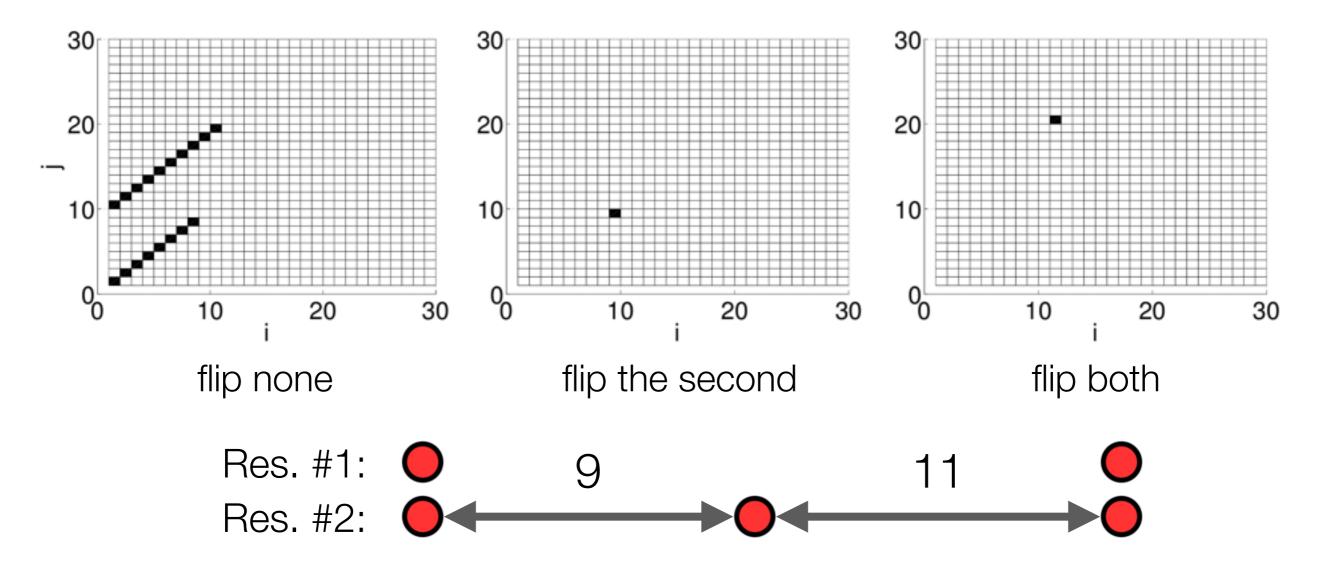
- We assume that intervals given by the strategy are
  - discrete (e.g., key or password renewal policy is defined in days or weeks) •
  - finite (i.e., every key or password is changed eventually) •
- $\rightarrow$  Markov strategy is defined by a finite set of probabilities
  - one for each subset of resources and each combination of times elapsed: (for example, with two resources,  $p^{S_{i,j}}$  is the probability of flipping subset S given that the first resource was flipped *i* steps ago and the second resource was flipped *j* steps ago)
- For a given strategy, we can find the optimal best-• response Markov strategy using linear programming
  - running time is exponential in the number of resources (



 on a desktop PC, easy for a few resources and dozens time intervals

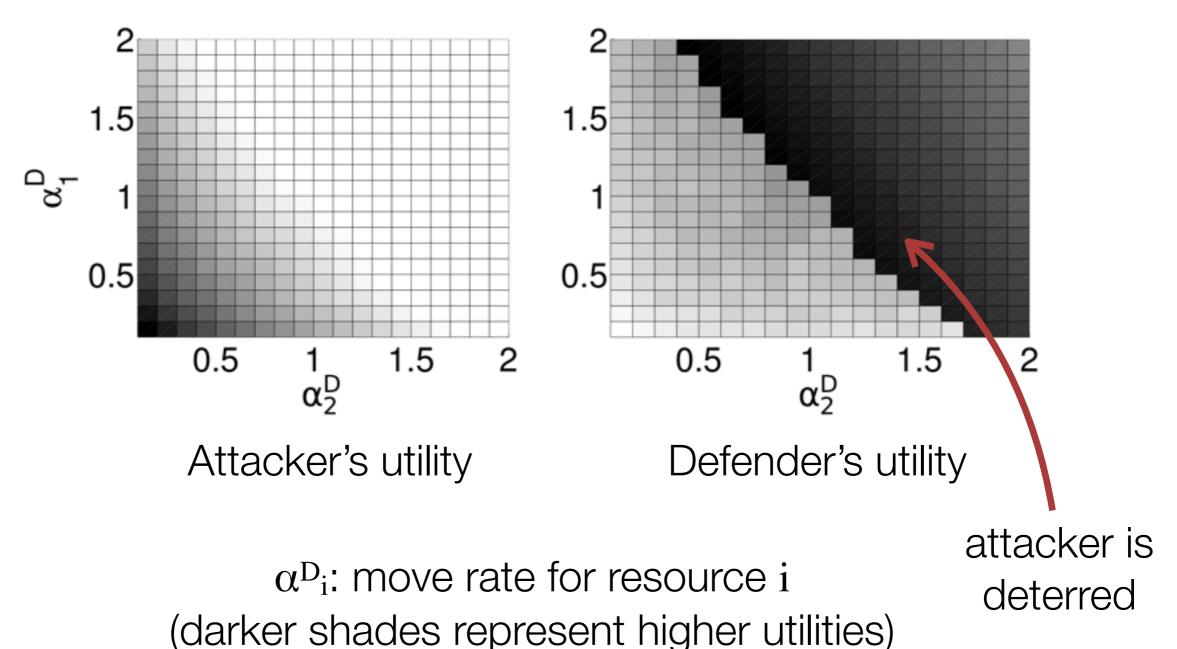
# Example: Markov Attack against a Given Defense

- Defender uses two independent exponential strategies with mean intervals 1 and 1/3
- Time steps are 0.03 long and the maximum number of time steps between two flips is 30



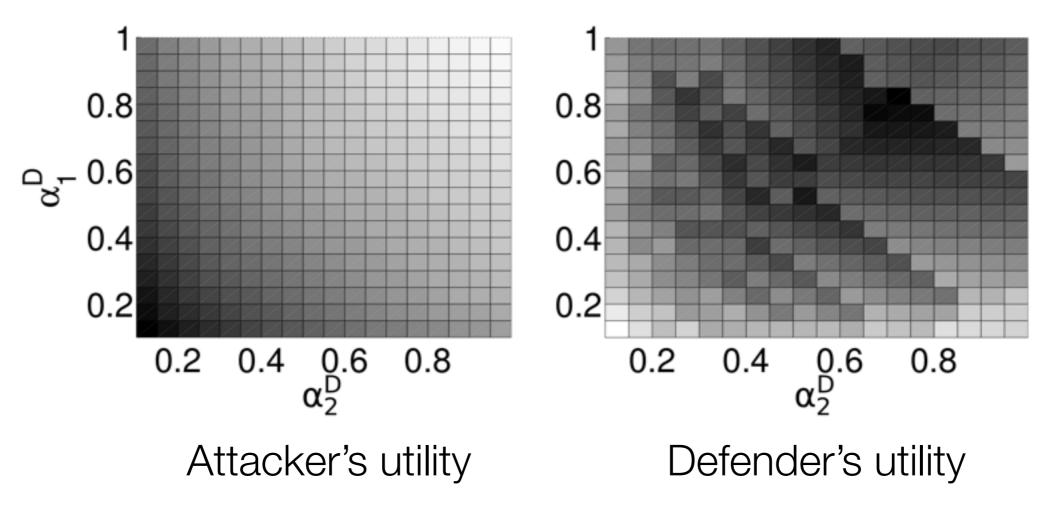
# Defense against a Markov Attacker (AND Model)

Defender uses independent periodic strategies



# Defense against a Markov Attacker (AND Model)

Defender uses independent exponential strategies



α<sup>D</sup><sub>i</sub>: move rate for resource i (darker shades represent higher utilities)

### Defense against a Markov Attacker - Lessons Learned

- Against a non-adaptive attacker, independent periodic strategies are good a choice in the AND model
  - however, an adaptive attacker could exploit this strategy
- Defender's utility is neither a continuous nor a monotonic function of the flipping rates, which makes optimization challenging
  - after the attacker has been deterred, increasing flipping rates only increases moving costs
  - with exponential strategies, the defender's utility has multiple local maxima

# Thank you for your attention!

Questions?

